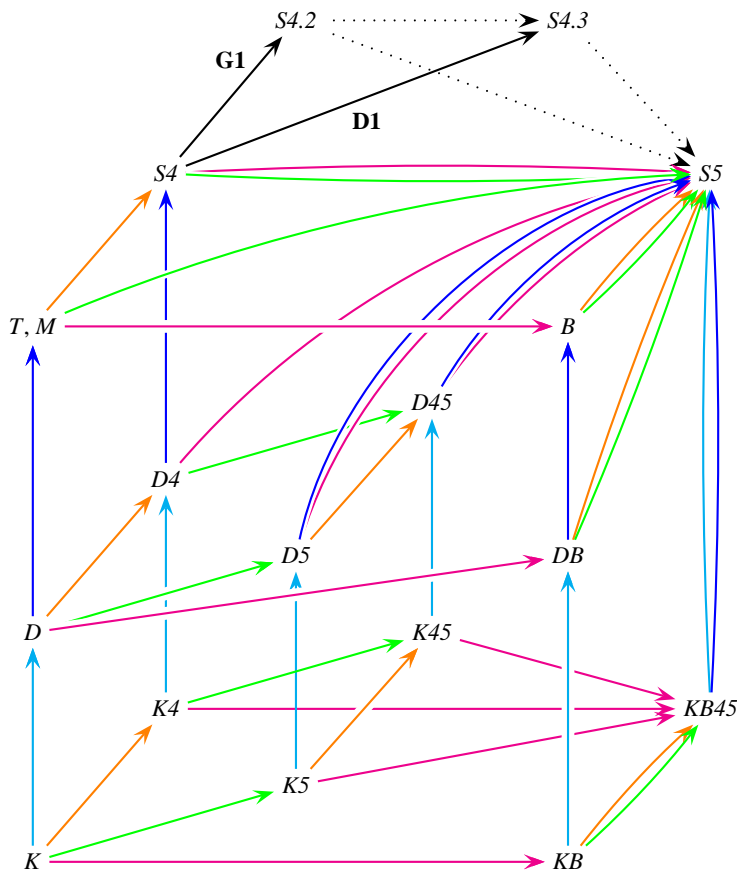


	if $\vdash \phi$ then $\vdash \Box \phi$			
N			Necessitation Rule	
K	$\Box(\phi \rightarrow \psi) \rightarrow \Box\phi \rightarrow \Box\psi$	$\Box(\phi \rightarrow \psi) \rightarrow \Diamond\phi \rightarrow \Diamond\psi$	Normality Axioms	
C	$\Box\phi \wedge \Box\psi \rightarrow \Box(\phi \wedge \psi)$	$\Diamond(\phi \vee \psi) \rightarrow \Diamond\phi \vee \Diamond\psi$	(by Normality and Necessitation)	
M	$\Box(\phi \wedge \psi) \rightarrow \Box\phi \wedge \Box\psi$	$\Diamond\phi \vee \Diamond\psi \rightarrow \Diamond(\phi \vee \psi)$	(by Normality and Necessitation)	
?	$\Box\phi \vee \Box\psi \rightarrow \Box(\phi \vee \psi)$	$\Diamond(\phi \wedge \psi) \rightarrow \Diamond\phi \wedge \Diamond\psi$	(by Normality and Necessitation)	
Scott–Lemmon	$\Diamond^h \Box^i \phi \rightarrow \Box^j \Diamond^k \phi$	$\Diamond^j \Box^k \phi \rightarrow \Box^h \Diamond^i \phi$	(h, i, j, k) -diamond property	$\forall w, x, y. wR^h x \wedge wR^j y \Rightarrow \exists z. xR^i z \wedge yR^k z$
T, M	$\Box\phi \rightarrow \phi$	$\phi \rightarrow \Diamond\phi$	reflexive	$\forall x. xRx$
B	$\phi \rightarrow \Box\Diamond\phi$	$\Diamond\Box\phi \rightarrow \phi$	symmetry	$\forall x, y. xRy \Rightarrow yRx$
B'	$\Box(\psi \vee \Box\phi) \rightarrow \phi \vee \Box\psi$	$\phi \wedge \Diamond\psi \rightarrow \Diamond(\psi \wedge \Diamond\phi)$	symmetry	$\forall x, y. xRy \Rightarrow yRx$
4	$\Box\phi \rightarrow \Box\Box\phi$	$\Diamond\Diamond\phi \rightarrow \Diamond\phi$	transitive	$\forall x, y, z. xRy \wedge yRz \Rightarrow xRz$
4_c	$\Box\Box\phi \rightarrow \Box\phi$	$\Diamond\phi \rightarrow \Diamond\Diamond\phi$	dense	$\forall x, z. xRz \Rightarrow \exists y. xRy \wedge yRz$
5	$\Diamond\phi \rightarrow \Box\Diamond\phi$	$\Diamond\Box\phi \rightarrow \Box\phi$	(right) Euclidean	$\forall x, y, z. xRy \wedge xRz \Rightarrow yRz$
D_c	$\Diamond\phi \rightarrow \Box\phi$	$\Diamond\phi \rightarrow \Box\phi$	functional	$\forall x, y, z. xRy \wedge xRz \Rightarrow y = z$
D	$\Box\phi \rightarrow \Diamond\phi$	<i>self-dual</i>	serial	$\forall x. \exists y. xRy$
D'	$\Box\neg\phi \rightarrow \neg\Box\phi$	<i>self-dual</i>	serial	$\forall x. \exists y. xRy$
G1, C	$\Diamond\Box\phi \rightarrow \Box\Diamond\phi$	<i>self-dual</i>	“convergent”; i.e., confluent	$\forall w, x, y. wRx \wedge wRy \Rightarrow \exists z. xRz \wedge yRz$
MS, M	$\Box\Diamond\phi \rightarrow \Diamond\Box\phi$	<i>self-dual</i>	?	(has no first-order frame condition)
$\Box T, \Box M, H$	$\Box(\Box\phi \rightarrow \phi)$		(right) shift-reflexive	$\forall x, y. xRy \Rightarrow yRy$
D1	$\Box(\Box\phi \rightarrow \psi) \vee \Box(\Box\psi \rightarrow \phi)$		“connected”; i.e., pre-Euclidean	$\forall x, y, z. xRy \wedge xRz \Rightarrow yRz \vee zRy$
GL, L, W	$\Box(\Box\phi \rightarrow \phi) \rightarrow \Box\phi$		transitive and well-founded	(has no first-order frame condition)



$$\Box(\pi) = \mathbf{K} \$ \mathbf{N}(\pi)$$

$$\mathbf{T} = \mathbf{B}^{op} \circ \mathbf{D} \circ \mathbf{4}$$

$$\mathbf{B} = \mathbf{5} \circ \mathbf{T}^{op}$$

$$\mathbf{4} = \Box(\mathbf{5}^{op}) \circ \mathbf{B}$$

$$\mathbf{5} = \Box(\mathbf{4}^{op}) \circ \mathbf{B}$$

$$\mathbf{D} = \mathbf{T}^{op} \circ \mathbf{T}$$

$$\mathbf{4} = \dots \mathbf{GL} \dots$$

$$\mathbf{4}_c = \mathbf{K} \$ \mathbf{H}$$

$$\mathbf{D1} = \dots \mathbf{5} \dots$$

Correspondence between N/K and \Box I/ \Box E

Let $_ _$ denote the modus ponens rule (associating to the left); let $\llbracket _ \rrbracket$ denote \Box I; and, let $\! _$ denote \Box E. For any context $C : \phi \vdash \psi$ and any proof $p : \Gamma \vdash \Box \phi$ we have the following functoriality theorem:

$$\begin{aligned} \llbracket _ \! _ \rrbracket : (\phi \vdash \psi) &\Rightarrow (\Gamma \vdash \Box \phi) \Rightarrow (\Gamma \vdash \Box \psi) \\ \llbracket C[!p] \rrbracket &= \mathbf{K} \ \$ \ \mathbf{N}(\lambda x. C[x]) \ \$ \ p \end{aligned}$$

If you have multiple proofs $p_i : \Gamma \vdash \Box \phi_i$, then just repeat the process as necessary; e.g.,

$$\llbracket C[!p, !q] \rrbracket = \mathbf{K} \ \$ \ \mathbf{K} \ \$ \ \mathbf{N}(\lambda x y. C[x, y]) \ \$ \ p \ \$ \ q$$

Thus, the \Box I rule corresponds to the **N** rule, and the \Box E rule corresponds to the **K** axiom.

Example proofs

K entails **C**

1	$\Box \phi \wedge \Box \psi$	
2	$\Box \phi$	\wedge E: 1
3	$\Box \psi$	\wedge E: 1
4	$\Box(\phi \rightarrow \psi \rightarrow \phi \wedge \psi)$	N
5	$\Box \phi \rightarrow \Box(\psi \rightarrow \phi \wedge \psi)$	\rightarrow E: K , 4
6	$\Box(\psi \rightarrow \phi \wedge \psi)$	\rightarrow E: 5, 2
7	$\Box \psi \rightarrow \Box(\phi \wedge \psi)$	\rightarrow E: K , 6
8	$\Box(\phi \wedge \psi)$	\rightarrow E: 7, 3
9	$\Box \phi \wedge \Box \psi \rightarrow \Box(\phi \wedge \psi)$	\rightarrow I: 1–8

1	$\Box \phi \wedge \Box \psi$	
2	$\Box \phi$	\wedge E: 1
3	$\Box \psi$	\wedge E: 1
4	* ϕ	\Box E: 2
5	ψ	\Box E: 3
6	$\phi \wedge \psi$	\wedge I: 4, 5
7	$\Box(\phi \wedge \psi)$	\Box I: 4–6
8	$\Box \phi \wedge \Box \psi \rightarrow \Box(\phi \wedge \psi)$	\rightarrow I: 1–7

K entails **M**

1	$\Box(\phi \wedge \psi)$	
2	* $\phi \wedge \psi$	\Box E: 1
3	ϕ	\wedge E: 2
4	$\Box \phi$	\Box I: 2–3
5	* $\phi \wedge \psi$	\Box E: 1
6	ψ	\wedge E: 5
7	$\Box \psi$	\Box I: 5–6
8	$\Box \phi \wedge \Box \psi$	\wedge I: 4, 7
9	$\Box(\phi \wedge \psi) \rightarrow \Box \phi \wedge \Box \psi$	\rightarrow I: 1–8

1	$\Box(\phi \wedge \psi)$	
2	$\Box(\phi \wedge \psi \rightarrow \phi)$	N
3	$\Box(\phi \wedge \psi) \rightarrow \Box \phi$	\rightarrow E: K , 2
4	$\Box \phi$	\rightarrow E: 3, 1
5	$\Box(\phi \wedge \psi \rightarrow \psi)$	N
6	$\Box(\phi \wedge \psi) \rightarrow \Box \psi$	\rightarrow E: K , 5
7	$\Box \psi$	\rightarrow E: 6, 1
8	$\Box \phi \wedge \Box \psi$	\wedge I: 4, 7
9	$\Box(\phi \wedge \psi) \rightarrow \Box \phi \wedge \Box \psi$	\rightarrow I: 1–8

Just a helpful proof to have around

1	□ϕ ∧ ◇ψ	
2	□¬(ϕ ∧ ψ)	
3	□ϕ	∧E: 1
4	* ϕ	□E: 3
5	¬(ϕ ∧ ψ)	□E: 2
6	¬ϕ ∨ ¬ψ	De Morgan: 5
7	¬ψ	Disjunctive Syllogism: 4, 6
8	□¬ψ	□I: 4–7
9	◇ψ	∧E: 1
10	¬□¬ψ	R: 9
11	⊥	¬E: 10, 8
12	¬□¬(ϕ ∧ ψ)	¬I: 2–11
13	◇(ϕ ∧ ψ)	R: 12
14	□ϕ ∧ ◇ψ → ◇(ϕ ∧ ψ)	→I: 1–13

GL entails 4

1	□ϕ	
2	* ϕ	□E: 1
3	□(ϕ ∧ □ϕ)	
4	□ϕ ∧ □□ϕ	→E: M , 3
5	□ϕ	∧E: 4
6	ϕ ∧ □ϕ	∧I: 2, 5
7	□(ϕ ∧ □ϕ) → ϕ ∧ □ϕ	→I: 3–6
8	□(□(ϕ ∧ □ϕ) → ϕ ∧ □ϕ)	□I: 2–7
9	□(ϕ ∧ □ϕ)	→E: GL , 8
10	□ϕ ∧ □□ϕ	→E: M , 9
11	□□ϕ	∧E: 10
12	□ϕ → □□ϕ	→I: 1–11