

	\mathcal{S}	\oplus	$\mathbb{0}$	\otimes	$\mathbb{1}$	\otimes - comm. \otimes - idem.	\oplus - idem.	$\overline{(\cdot)}$	\ominus	$(\cdot)^*$
natural numbers	\mathbb{N}	$+$	0	$*$	1	\checkmark \times	\times		\times	\times
compactification of \mathbb{N}	$\mathbb{N} \cup \{+\infty\}$	$+$	0	$*$ ¹	1	\checkmark \times	\times		\times	$(x \equiv 0) ? 1 : +\infty$
non-negative rationals	$\mathbb{Q}_{[0,+\infty)}$	$+$	0	$*$	1	\checkmark \times	\times		\times	\times
non-negative reals	$\mathbb{R}_{[0,+\infty)}$	$+$	0	$*$	1	\checkmark \times	\times		\times	\times
log-domain non-negative reals	$\mathbb{R}_{[-\infty,+\infty)}$	log sum	$-\infty$	$+$	0	\checkmark \times	\times		\times	\times
integers	\mathbb{Z}	$+$	0	$*$	1	\checkmark \times	\times	$-x$	$x-y$	\times
integers modulo n	$\mathbb{Z}/n\mathbb{Z}$	$+$	0	$*$	1	\checkmark \times	\times	$-x$	$x-y$	\times
rationals	\mathbb{Q}	$+$	0	$*$	1	\checkmark \times	\times	$-x$	$x-y$	\times
reals	\mathbb{R}	$+$	0	$*$	1	\checkmark \times	\times	$-x$	$x-y$	\times
one-point compactification ² of \mathbb{R}	$\mathbb{R} \cup \{\infty\}$	$+$ ³	0	$*$ ¹	1	\checkmark \times	\times	\times ⁴		$1/(1-x)$ ⁵
extended Viterbi	$\mathbb{R}_{[0,+\infty]}$	max	0	$*$ ¹	1	\checkmark \times	\checkmark		\times	$(x \leq 1) ? 1 : +\infty$
Viterbi	$\mathbb{R}_{[0,+\infty)}$	max	0	$*$	1	\checkmark \times	\checkmark		\times	\times
log-domain Viterbi	$\mathbb{R}_{[-\infty,+\infty)}$	max	$-\infty$	$+$	0	\checkmark \times	\checkmark		\times	\times
probabilistic Viterbi	$\mathbb{R}_{[0,1]}$	max	0	$*$	1	\checkmark \times	\checkmark		\times	1
log-domain probabilistic Viterbi	$\mathbb{R}_{[-\infty,0]}$	max	$-\infty$	$+$	0	\checkmark \times	\checkmark		\times	0
non-negative tropical	$\mathbb{R}_{[0,+\infty]}$	min	$+\infty$	$+$	0	\checkmark \times	\checkmark		\times	0
tropical	$\mathbb{R}_{(-\infty,+\infty]}$	min	$+\infty$	$+$	0	\checkmark \times	\checkmark		\times	\times
extended tropical	$\mathbb{R}_{[-\infty,+\infty]}$	min	$+\infty$	$+$ ⁶	0	\checkmark \times	\checkmark		\times	$(x \geq 0) ? 0 : -\infty$
bounded distributive lattice, L	L	\sqcup	\perp	\sqcap	\top	\checkmark \checkmark	\checkmark		\times	\top
bounded totally ordered set, X	$X_{[a..b]}$	max	a	min	b	\checkmark \checkmark	\checkmark		\times	b
binary Boolean algebra	$\{0, 1\}$	\vee	0	\wedge	1	\checkmark \checkmark	\checkmark		\times	1
subsets of X , as a Boolean algebra	$\mathcal{P}(X)$	\cup	\emptyset	\cap	X	\checkmark \checkmark	\checkmark		\times	X
binary Boolean ring	$\{0, 1\}$	$\underline{\vee}$ ⁷	0	\wedge	1	\checkmark \checkmark	\times	x ⁷	$x \underline{\vee} y$ ⁷	\times
subsets of X , as a Boolean ring	$\mathcal{P}(X)$	Δ ⁸	\emptyset	\cap	X	\checkmark \checkmark	\times	x ⁸	$x \Delta y$ ⁸	\times
regular languages over Σ	$\mathcal{P}(\Sigma^*)$	\cup	\emptyset	\cdot	$\{\varepsilon\}$	\times \times	\checkmark		\times	x^*
square matrices over semiring R	$R(n, n)$	$+$	$\mathbf{0}_n$	$*$	\mathbf{I}_n	\times \times	\checkmark ⁹	$-x$ ¹⁰	$x - y$ ¹⁰	x^{*11}

¹ Where we define $0 * +\infty = 0 = +\infty * 0$ or, as appropriate, $0 * \infty = 0 = \infty * 0$.

² Note that our use of the Alexandroff compactification is distinct from the real projective line ($\hat{\mathbb{R}}$, $\mathbb{R}P^1$, or $P^1(\mathbb{R})$). The real projective line interpretation of $\mathbb{R} \cup \{\infty\}$ requires leaving the following six expressions undefined: $\infty + \infty$, $\infty - \infty$, $0 * \infty$, $\infty * 0$, ∞ / ∞ , and $0/0$. However, to properly form a $*$ -semiring we require defining all but the last two.

³ Where we define $\infty + \infty = \infty$.

⁴ Beware, this is not a ring! The natural negation operation has that $-\infty = \infty$, so negation does not always give the additive inverse. Together with our definition of addition this means: $\infty - \infty = \infty + -\infty = \infty + \infty = \infty$.

⁵ Except that we define $1^* = \infty$ and $\infty^* = \infty$.

⁶ Where we define $+\infty + -\infty = +\infty$.

⁷ That is, exclusive or: $x \underline{\vee} y = (x \vee y) \wedge \neg(x \wedge y) = (x \wedge \neg y) \vee (y \wedge \neg x) = \neg(x \Leftrightarrow y)$. Note that $x \underline{\vee} x = 0$. Thus, do not confuse the Boolean ring's negation (the identity function) with the Boolean algebra's complementation (i.e., \neg).

⁸ That is, symmetric difference: $x \Delta y = (x \cup y) \setminus (x \cap y) = (x \setminus y) \cup (y \setminus x)$. Note that $x \Delta x = \emptyset$. Thus, do not confuse the Boolean ring's negation (the identity function) with the Boolean algebra's complementation (i.e., $\bar{\cdot}$).

⁹ Provided R forms a dioid (i.e. \oplus -idempotent semiring).

¹⁰ Provided R forms a ring.

¹¹ Provided R forms a $*$ -semiring (i.e., closed semiring). An efficient implementation is the famous Gauss–Jordan–Floyd–Warshall–McNaughton–Yamada algorithm. Beware that the notation for asteration should not be confused with the antiautomorphic involution (e.g., conjugate transpose) of \star -rings (i.e., ring with involution).